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Der Präsident des Europäischen Patentamts;
Im Auftrag

For the President of the European Patent Office

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Blatt 2 der Bescheinigung
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CYCLE SYNCHRONIZATION BETWEEN INTERCONNECTED SUB-NETWORKS

Description

1 The present invention relates to a method to perform a cycle synchronization between interconnected sub-networks and a cycle synchronizator adapted to perform said method.

5 It is known to interconnect sub-networks, e. g. with long delay bi-directional connections to extend a network to a wider area. In particular, this technique is used to interconnect several IEEE 1394 serial buses to extend an IEEE 1394 network, e. g. through a whole house. The basic topology of such a connection is shown in Fig. 1. A first interface 20 is part of a first IEEE 1394 serial bus
10 21 which might consist of a number of IEEE 1394 nodes. A second interface 22 is part of a second IEEE 1394 serial bus 23 which might comprise another number of IEEE 1394 nodes. The first interface 21 and the second interface 22 are connected via a long delay bi-directional connection 24 which might be, but is not restricted to a coax cable medium.

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Independent IEEE 1394 buses must be synchronized to have the same cycle rate. In particular, the IEEE 1394 standard mandates that for opened isochronous channels an isochronous packet is sent in every isochronous cycle. To ensure that isochronous transfers between the interconnected IEEE 1394
20 buses work, it must be ensured that all buses have the same frequency of isochronous cycles.

Therefore, it is the object underlying the present invention to provide a method to perform a cycle synchronization between interconnected sub-networks and a
25 cycle synchronizator adapted to perform said method.

The method according to the present invention is defined in independent claim 1 and the cycle synchronizator according to present invention is defined in independent claim 12. Preferred embodiments thereof are respectively defined
30 in the dependent subclaims.

The method to perform a cycle synchronization between interconnected sub-networks according to the present invention is characterized in that a reference node connected to one of the sub-networks transmits a respective
35 cycle time information to cycle masters of all other sub-networks at recurring

- 1 time instants, and the cycle masters of all other sub-networks adjust their cycle time accordingly.

Therewith, the present invention offers a method to synchronize several inter-
5 connected sub-networks which is independent of the connection between the sub-networks, since with the transmission of cycle time information of a reference node no relying on a clock frequency used for the transmission through the connection through sub-networks is necessary. After reception of the cycle time information each cycle master of the other sub-networks can
10 adjust their cycle time accordingly so that in turn the cycle frequency in an IEEE 1394 serial bus connected to a respective cycle gets adjusted. Therefore, in a network with N sub-networks N-1 cycle masters are required to adjust their cycle time and the remaining sub-network has to comprise the reference node transmitting its time information to the N-1 cycle masters of the other
15 sub-networks. Preferably, the reference node and the cycle masters are arranged within a respective interface of the sub-network which is connected to the interconnection of all sub-networks.

According to the present invention an adjustment of the cycle time within a
20 cycle master might be performed by the following steps: Determining a first time interval in-between two receptions of cycle time information from the reference node with an own clock, determining a second time interval in-between corresponding transmission of cycle time information from the reference node on basis of the received cycle time information, comparing the
25 first and second time intervals and adjusting the own cycle time according to the comparison result. Therefore, a large scale integration is possible.

Further, the comparison of the first and the second time intervals according to the present invention might consider a preceding adjustment of the own cycle
30 time, the adjustment of the own cycle time within a cycle master might be performed in a step-wise manner and/or the adjustment of the own cycle time within a cycle master might be performed by adjusting a local number of clocks within one cycle.

35 In particular, in the latter case the adjustment of the own cycle time within a cycle master is performed by setting the local number of clocks equal to an ideal number of clocks of one cycle in case the first time interval and the second time interval are identical, smaller than an ideal number of clocks of

- 1 one cycle in case the first time interval is smaller than the second time interval
and larger than an ideal number of clocks in case the first time interval is
larger than the second time interval. In particular, these features enable a very
easy and therefore reliable method to perform the cycle synchronization
5 between interconnected sub-networks according to the present invention which
is independent from the transmission method used in-between the sub-net-
works.

The stepwidth of setting the local number of clocks smaller or larger than the
10 ideal number of clocks might be determined according to the difference of the
first and second time intervals. In this case it is possible to determine how fast
the sychronization should be achieved and/or to consider smaller and larger
deviations of the cycle timers within the cycle masters.

- 15 According to the present invention preferably the cycle time information trans-
mitted by the reference node is a content of its cycle time register. In this case
the adjustment of the own cycle time within a cycle master is preferably
performed by adjusting the average difference between a time interval of two
transmissions of cycle time information of the reference node which is deter-
20 mined by subtracting two succeeding received contents of the cycle time
register of the reference node and a time interval of two samplings of the own
cycle timer which is determined by subtracting two succeeding sampled
contents of the own cycle time register plus a corrective difference to be zero.
Of course, also other than two succeeding transmissions could be used, but in
25 this case the hardware design to realize a cycle synchronizator according to
the present invention leads to an increased cost. Further preferably, the cor-
rective difference corresponds to the preceding adjustment.

- Further preferably, according to the present invention the recurring time
30 instants are determined to a regular time interval with a small variation.

- The cycle synchronizator according to the present invention is therefore
characterized by a clock offset estimation means to determine a timing error of
an own cycle timer, and a cycle adjustment loop receiving the timing error
35 determined by said clock offset estimation means to adjust the own cycle timer
to reduce its timing error. Preferably, a de-jitter filter is arranged in-between
the clock offset estimation means and the cycle adjustment loop to filter said
determined timing error.

1 Therefore, in case the present invention is applied to a distributed IEEE 1394
network, i. e. several IEEE 1394 serial buses which are regarded as sub-net-
works are interconnected, e. g. by a long delay, bi-directional connection
provides advantages in that the cycle synchronization is based on free-running
5 oscillators of the cycle masters and standard IEEE 1394 physical interfaces
can be used, since the cycle synchronization is based on a timing error of the
own cycle timer which can be determined on basis of the transmission of cycle
time information of a reference node in the network. Additionally, the reference
node does not need to be a cycle master, i. e. the reference node can be prede-
10 terminated.

Further features and advantages of the present invention will be apparent from
the following detailed description of an exemplary embodiment thereof taken in
conjunction with the accompanying drawings, in which

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Fig. 1 shows an overview of a simple long delay IEEE 1394 network,

Fig. 2 shows a timing diagram of a first preferred embodiment
according to the present invention, and

20

Fig. 3 shows a phase locked loop for cycle synchronization according
to a preferred embodiment of the present invention.

The following preferred embodiment of the present invention is adapted to the
25 IEEE 1394 standard. However, as mentioned above, the present invention is
not restricted thereto.

Every IEEE 1394 node maintains cycle time information. This is basically a
register that is incremented by a local, free-running clock of 24.576 MHz or in-
30 teger multiples of that. According to the present invention this cycle time
information is transmitted at regular instants via the interconnection of
several sub-networks, in case of the example shown in Fig. 1 via the long delay
bi-directional connection 24. The basic assumption of this method is that
transmission of the cycle occurs at recurring time instants, preferably regular
35 intervals, e. g. every 10 ms. Further, the exact value of that interval is not im-
portant since the exact value can be recovered from the difference of two
transmitted samples of the cycle time register and the corresponding time

- 1 stamps of the receiver will be sampled at the instant when the transmitted samples are received.

Fig. 2 shows an example of timing of the transmission and reception of the cycle time. The node which has been chosen to be the reference node transmits the time at least to all other nodes comprising a cycle master. As mentioned above, the reference node is not required to be the cycle master within its connected IEEE 1394 sub-network. As shown in Fig. 2, the reference node samples its local cycle time register at regular instants, i. e. at a first transmission time t_0 , a second transmission time t_3 , and a third transmission time t_5 at which the contents of the cycle time register are respectively transmitted. It is also shown in Fig. 2 that the second transmission time t_3 , which is an actual transmission time, differs from an ideal second transmission time t_2 by a time difference $t_{\text{jitter}1}$. After a respective transmission of the contents of the cycle time register these contents are received at a first reception time t_1 , a second reception time t_4 , and a third reception time t_6 . Similar to the case of the transmission it is shown in Fig. 2 that the actual reception of the transmitted cycle time register content at the second reception time t_4 differs from an ideal reception thereof. The difference in-between the ideal and the later actual second reception time is labelled with $t_{\text{jitter}2}$. A difference in-between the first and second actual transmission times is determined to Δt_2 and in-between the second and third actual transmission times to $\Delta t_2'$. A difference in-between the first and second actual reception times is determined to Δt_1 and in-between the second and third actual reception times to $\Delta t_1'$.

25 To allow for significant jitter to occur both on the transmitter and on the receiver side according to the present invention an optional filtering can be performed which limits the cycle length adjustment range to ± 1 clock and/or which uses a de-jitter filter.

30 After transmission, the receiving node samples its own local cycle timer at the instant when it receives the remote cycle time information. In the standard IEEE 1394 node, one cycle has a duration of 3072 clocks of a 24,576 MHz oscillator. According to the preferred embodiment of the present invention shown and described in the following, a cycle timer is used where the duration of the cycle can be adjusted to 3071, 3072 or 3073 clocks. However, a variable duration might also be implemented. The information of the remote and local cycle time registers is used to adjust the local number of clocks per cycle.

- 1 According to the preferred embodiment of the present invention a special phase locked loop as shown in Fig. 3 is used to achieve the synchronization.

The cycle synchronizator shown in Fig. 3 comprises a clock offset estimation means to determine a timing error in clocks which is supplied to a cycle adjustment loop 2 preferably, as shown in Fig. 3, via a de-jitter filter 4. The cycle adjustment loop 2 in turn determines a new cycle duration which is supplied back to the clock offset estimation means 1.

- 10 In particular, the clock offset estimation means 1 receives the remote time which is supplied directly as minuend to a first adder 9 and via a first delay element 11 as subtrahend to the first adder 9. The first delay element 11 holds the preceding sample of the remote time, i. e. shows a FIFO-behaviour with a storage capacity of one sample. Therefore, the first adder 9 outputs a remote time delta, i. e. the time difference of the time in-between two samples of the time register of the reference node. This remote time delta is input as minuend to a second adder 5. Further, the clock offset estimation means 1 comprises the local cycle timer 3 of the cycle master. The local time output therefrom is input as minuend to a third adder 10 and via a second delay element 12 also as subtrahend to the third adder 10. The second delay element 12 shows the same delay T as the first delay element 11. Therefore, the third adder 10 outputs a local time delta corresponding in time to the remote time delta output by the first adder 9. This local time delta is input as subtrahend to the second adder 5 which outputs the timing error in clocks to the de-jitter filter 4 which inputs the filtered timing error in clocks to the cycle adjustment loop 2. Further, the local time delta output from the third adder 10 is input to a controller 7 which determines the number clock skips/inserts necessary on basis of an arithmetic operation subtracting the duration of a cycle in clocks for this period of time from that of an ideal cycle and multiplying the resulting difference by the quotient of the number of clocks between the previous sampling instant and this sampling instant with the duration of a cycle in clocks for this period of time.

- The cycle adjustment loop 2 comprises a fourth adder receiving the timing error in clocks from de-jitter filter 4 as a first summand and the number of clock skips/inserts determined by the controller 7 as second summand to build their sum. This sum is supplied to an integrator 13 which outputs its integration result to a quantizer 6. The quantizer 6 determines whether the next

1 cycle of the cycle timer 3 within the clock offset estimation means 1 should
comprise 3071, 3072 or 3073 clocks. In case the integration result of the inte-
grator 13 is smaller than -80 then the next cycle should comprise 3071
clocks, in case the integration result is bigger than 80 then the next cycle
5 should comprise 3073 clocks and in case the output result of the integrator 13
equals to 80 the cycle should comprise 3072 clocks. This comparison intro-
duces an hysteresis into the loop so that there are usually only differences of
one clock in succeeding cycles, i. e. that there is usually no jump from 3071 to
3073, but either between 3072 and 3073 or between 3071 and 3072 clocks per
10 cycle. Therefore, also another value than 80 cycles which equal to 10 ms might
be used. The number of clocks output by the quantizer 6 is input to a third de-
lay element 14 which also has the same delay T as the first delay element 11.
The cycle duration output by the third delay element 14 is supplied to the con-
troller 7 which determines the number of clock skips/inserts and to the cycle
15 timer 13.

As mentioned above, the delay T of the first to third delay elements 11, 12, 14
are not fixed but depend on the reception of a transmitted remote time. Also,
the delay T within the delay elements does not indicate a fixed or preset time,
20 but that the sample and hold operation performed by the delay element is
performed by all three delay elements simultaneously.

The phase locked loop for a cycle synchronization according to the preferred
embodiment of the present invention shown in Fig. 3 adjusts the average
25 difference between the remote time interval which is measured with the remote
clock and the local time interval which is measured with the local clock plus a
corrective difference to be zero.

Since without jitter or disturbances the delay of the transmission path
30 between reference and cycle synchronizator is constant, the method according
to the present invention uses exactly the time interval for the local and remote
measurement. Since the respective oscillators used for the respective measure-
ment might differ slightly with respect to their oscillation frequency, i. e.
according to the IEEE 1394 standard +/- 100 ppm are allowed, these measure-
35 ments of local and remote time intervals do not give exactly the same number
of clocks. The cycle synchronizator according to the present invention extracts
the number of cycles n_cycles that have elapsed in the respective time interval
and depending on the current cycle duration, the corrective number of clocks

1 is set to be either +n_cycles, 0. or -n_cycles. Corrective values of -1,0 + 1 per
3072 clocks as explained above is equivalent to an adjustment range of
+/- 166/3072 = +/- 325 ppm. Also larger corrective values might be used
which - on the other hand - lead to higher local jitter and are therefore not
5 preferred. Therefore, in the long run, the remote and local number of cycles
are equalized.

As shown in Fig. 3 it is advantageous to insert a de-jitter filter 4 before the
cycle timer adjustment loop. A suitable filter is a lowpass filter, but other
10 filters, e. g. a running mean or a time-adaptive lowpass may also be used. By
choosing a suitably high time constant in that filter which is independent from
the clock synchronization loop the jitter can be eliminated.

Since the IEEE 1394 serial bus is a self-configuring bus, it is necessary that
15 the network reference node is automatically determined after each network re-
configuration, e. g. after the addition or removal of nodes.

Therefore, according to the present invention the oscillator is not adjusted, but
the number of clocks within one cycle. Therefore, a free-running oscillator can
20 be used instead of a voltage controlled oscillator. This feature enables the
integration of the cycle synchronization according to the present invention on
a single chip. Further, as mentioned above, the present invention performs the
cycle synchronization independently of the connecting channel in-between the
different sub-networks, i. e. IEEE 1394 serial buses which basically need only
25 slight modifications in that the cycle synchronizator according to the present
invention has to be included into a respective cycle master of a sub-network.
Further, the connection network needs no master clock, since one of the sub-
networks serves as reference.

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Claims

- 1 1. Method to perform a cycle synchronization between interconnected sub-networks, **characterized in that**
- a reference node connected to one of the sub-networks transmits a respective cycle time information to cycle masters of all other sub-networks at
 - 5 recurring time instants, and
 - the cycle masters of all other sub-networks adjust their cycle time accordingly.
2. Method according to claim 1, **characterized in that** an adjustment of
- 10 the cycle time within a cycle master is performed by the following steps:
- determining a first time interval (Δt_1 , $\Delta t_1'$) in-between two receptions of cycle time information from the reference node with an own clock,
 - determining a second time interval (Δt_2 , $\Delta t_2'$) in-between two corresponding transmissions of cycle time information from the reference node on
 - 15 basis of the received cycle time information,
 - comparing the first time interval (Δt_1 , $\Delta t_1'$) and the second time interval (Δt_2 , $\Delta t_2'$), and
 - adjusting the own cycle length according to the comparison result.
- 20 3. Method according to claim 2, **characterized in that** the comparison of the first time interval (Δt_1 , $\Delta t_1'$) and the second time interval (Δt_2 , $\Delta t_2'$) considers a preceding adjustment of the own cycle length.
4. Method according to claim 2 or 3, **characterized in that** the adjustment
- 25 of the own cycle length within a cycle master is performed in a step-wise manner.
5. Method according to claim 2, 3 or 4, **characterized in that** the adjustment of the own cycle length within a cycle master is performed by adjusting a
- 30 local number of clocks within one cycle.
6. Method according to claim 5, **characterized in that** the adjustment of the own cycle length within a cycle master is performed by setting the local number of clocks
- equal to an ideal number of clocks of one cycle in case the first time

- 1 interval (Δt_1 , $\Delta t_1'$) and the second time interval (Δt_2 , $\Delta t_2'$) are identical,
- smaller than an ideal number of clocks of one cycle in case the first
time interval (Δt_1 , $\Delta t_1'$) is smaller than the second time interval (Δt_2 , $\Delta t_2'$), and
- larger than an ideal number of clocks in case the first time interval
5 (Δt_1 , $\Delta t_1'$) is larger than the second time interval (Δt_2 , $\Delta t_2'$).

7. Method according to claim 6, **characterized in that** a step-width to
adjust the own cycle timer within a cycle master is set according to the differ-
ence of the first time interval (Δt_1 , $\Delta t_1'$) and the second time interval (Δt_2 ,
10 $\Delta t_2'$).

8. Method according to anyone of the preceding claims, **characterized in
that** the cycle time information transmitted by the reference node is a content
of its cycle time register.
15

9. Method according to claim 8, **characterized in that** the adjustment of
the own cycle time within a cycle master is performed by adjusting the average
difference between a time interval of two transmissions of cycle time informa-
tion of the reference node which is determined by subtracting two succeeding
20 received contents of the cycle time register of the reference node and a time in-
terval of two samplings of the own cycle timer which is determined by sub-
tracting two succeeding sampled contents of the own cycle time register plus a
corrective difference to be zero.

25 10. Method according to claim 9, **characterized in that** the corrective differ-
ence corresponds to the preceding adjustment.

11. Method according to anyone of the preceding claims, **characterized in
that** the recurring time instants are determined according to a regular time in-
30 terval with a small variation.

12. Cycle synchronizator, **characterized by**
- a clock offset estimation means (1) to determine a timing error of an
own cycle timer (3), and
35 - a cycle adjustment loop (2) receiving the timing error determined by
said clock offset estimation means (1) to adjust the own cycle timer (3) to re-
duce its timing error.

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Sony International (Europe) GmbH

File: 52.652

10.10.2000

- 1 13. Cycle synchronizator according to claim 12, **characterized by** a de-jitter filter (4) arranged in-between the clock offset estimation means (1) and the cycle adjustment loop (2) to filter said determined timing error.

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Abstract**CYCLE SYNCHRONIZATION BETWEEN INTERCONNECTED SUB-NETWORKS**

A method to perform a cycle synchronization between several interconnected sub-networks, comprises the steps that a reference node connected to one of the sub-networks transmits a respective cycle time information to cycle masters of all other sub-networks at recurring time instants, and the cycle masters of all other sub-networks adjust their cycle time accordingly. Therefore, a cycle synchronizator comprises a clock offset estimation means (1) to determine a timing error of an own cycle timer (3), and a cycle adjustment loop (2) receiving the timing error determined by said clock offset estimation means (1) to adjust the own cycle timer (3) to reduce its timing error.

(Fig. 3)

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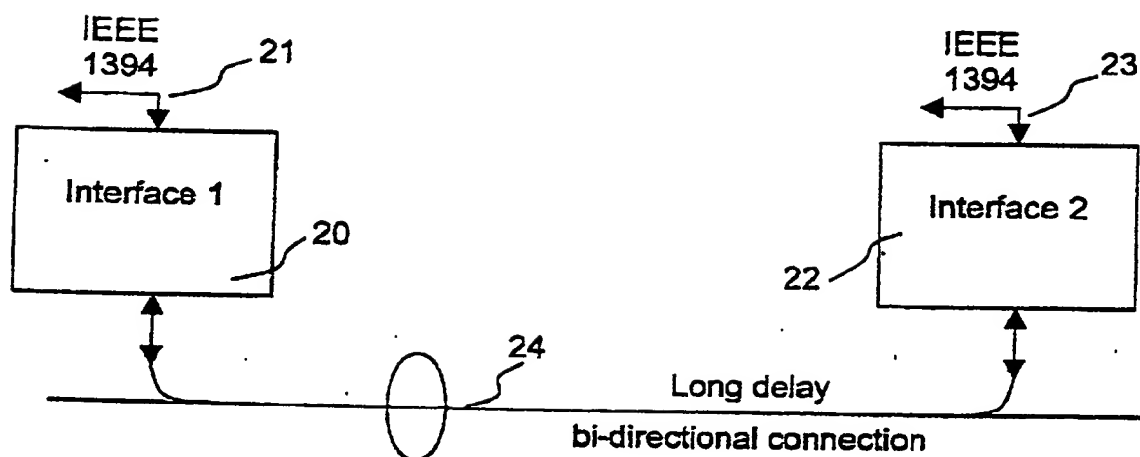
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10. Okt. 2000

Figure 1

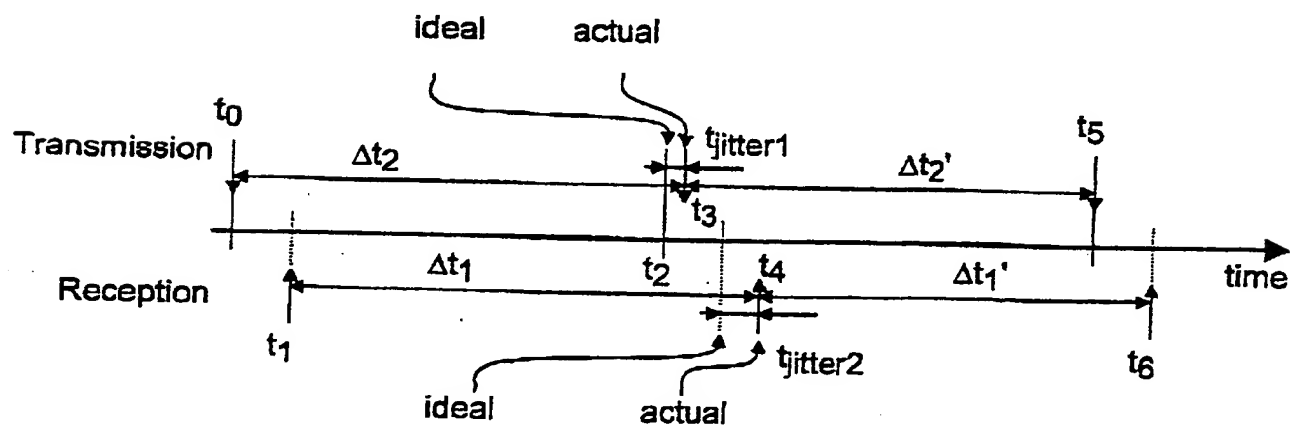


Figure 2

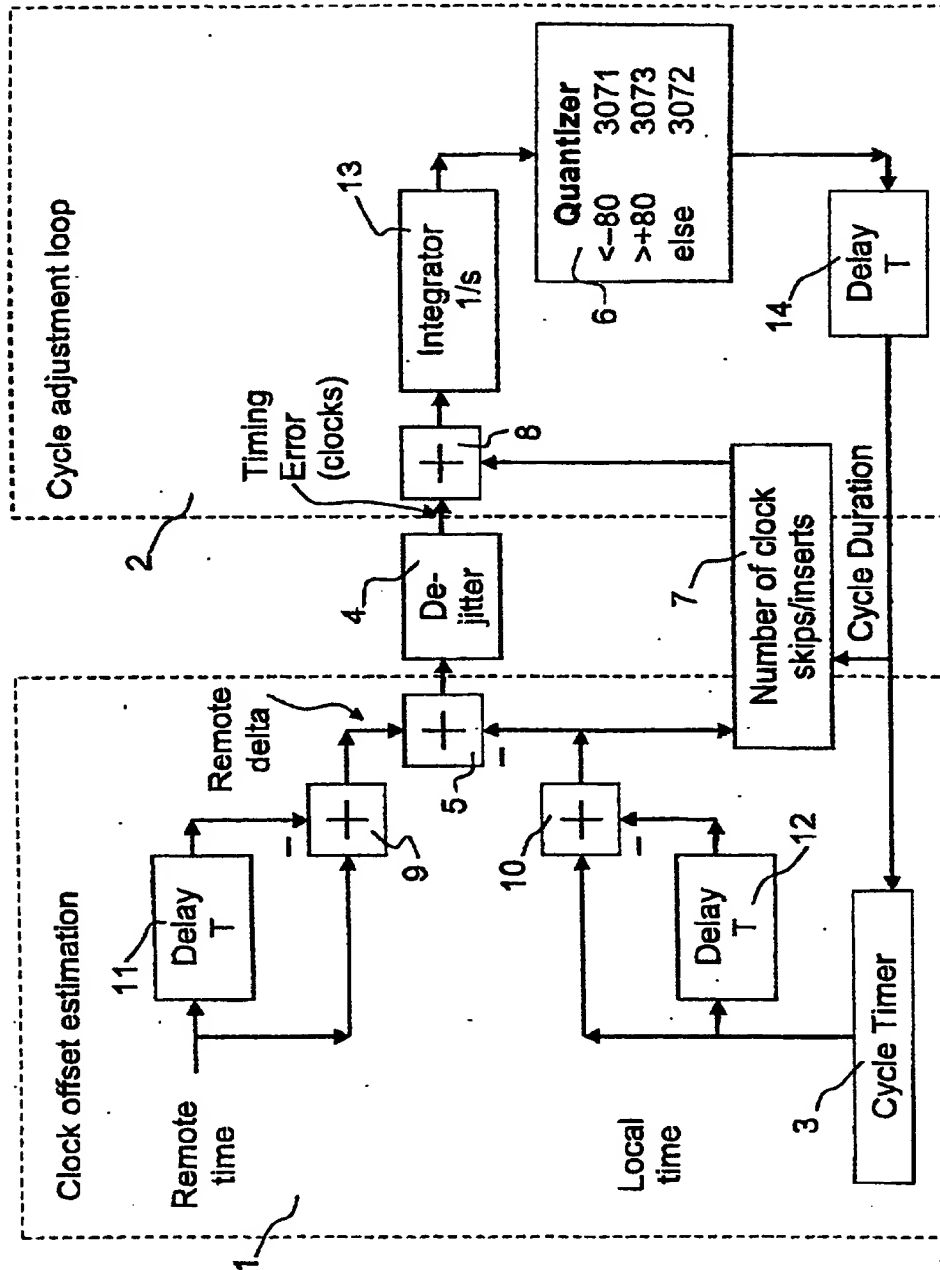


Figure 3